# Random sampling for the monomer-dimer model on a lattice

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In the monomer-dimer model on a graph, each matching (collection of nonover-lapping edges) M has a probability proportional to  $\lambda^{|M|}$ , where  $\lambda > 0$  is the model parameter, and |M| denotes the number of edges in M. An approximate random sample from the monomer-dimer distribution can be obtained by running an appropriate Markov chain (each step of which involves an elementary local change in the configuration) sufficiently long. Jerrum and Sinclair have shown (roughly speaking) that for an arbitrary graph and fixed  $\lambda$  and  $\epsilon$  (the maximal allowed variational distance from the desired distribution),  $O(|\Lambda|^2|E|)$  steps suffice, where |E| is the number of edges and  $|\Lambda|$  the number of vertices of the graph. For sufficiently nice subgraphs (e.g., cubes) of the d-dimensional cubic lattice we give an explicit recipe to generate approximate random samples in (asymptotically) significantly fewer steps, namely (for fixed  $\lambda$  and  $\epsilon$ )  $O(|\Lambda|(\ln|\Lambda|)^2)$ . © 2000 American Institute of Physics. [S0022-2488(00)01403-1]

#### I. INTRODUCTION

The monomer-dimer model, described below, originates from Statistical Physics, where it has been used to study the absorption of oxygen molecules on a surface, and the properties of a binary mixture. See Heilmann and Lieb<sup>8</sup> for further background and references. More recently, the model has also drawn much attention in Operations Research, Combinatorics, and Graph Theory (see Refs. 9 and 10).

Throughout this paper, the size (number of elements) of a finite set A will be denoted by |A|. Consider a finite, undirected graph  $G = (\Lambda, E)$ , where  $\Lambda$  is the set of vertices of G and E is the set of edges. A *matching* on G is a subset  $M \subset E$  such that no two edges in M have a common end point. Let  $\lambda > 0$  (this is the model parameter). Now assign to each matching M a probability proportional to  $\lambda^{|M|}$ .

Alternatively, define the state space  $\Omega = \{0,1\}^E$ . Elements of  $\Omega$  (called *configurations* on E) are typically denoted by  $\omega = (\omega_e, e \in E)$ . The monomer-dimer distribution for G (with parameter  $\lambda$ ) is then defined as

$$\mu(\omega) = \frac{\lambda^{|\omega|} I(\omega \text{ is allowable})}{Z(\lambda)},$$
(1.1)

where " $\omega$  is allowable" means that the set  $\{e:\omega_e=1\}$  is a matching,  $|\omega|$  denotes the size of that set, and  $Z(\lambda)$  is the normalization factor (partition function). It is clear that the two descriptions (one with state space the set of all matchings, the other with state space  $\{0,1\}^E$ ), are equivalent, and both descriptions will be used in this paper.

To continue, we need some more terminology and notation:

If two vertices i and j are adjacent, we write  $i \sim j$ . The  $degree \ deg(v)$  of a vertex v is defined as the number of edges that have v as an end point. If two edges  $e_1$  and  $e_2$  share a common end point, we write  $e_1 \sim e_2$ . Let, in the rest of this subsection,  $\Delta$  be a subset of E. We denote the set  $\{0,1\}^{\Delta}$  by  $\Omega_{\Delta}$ . Similarly, if  $\omega \in \Omega$ , then  $\omega_{\Delta}$  denotes the "restriction" of  $\omega$  to  $\Delta$ , i.e., the element  $(\omega_e : e \in \Delta)$  of  $\Omega_{\Delta}$ . If  $\omega$ ,  $\omega' \in \Omega_{\Delta}$  we call an edge  $e \in \Delta$  an edge of disagreement (w.r.t. the pair

 $(\omega, \omega')$ ) if  $\omega_e \neq \omega'_e$ , and we denote the set of all such edges by  $V(\omega, \omega')$ . The *boundary* of  $\Delta$ , denoted by  $\partial \Delta$ , consists of all elements  $e \in E \setminus \Delta$  such that  $e \sim e'$  for some  $e' \in \Delta$ .

Let  $\alpha \in \Omega_{\partial \Delta}$ , (so  $\alpha$  is a configuration on the boundary of  $\Delta$ ). The monomer-dimer distribution for  $\Delta$  with boundary condition  $\alpha$  is defined as follows:

$$\mu_{\Delta}^{\alpha}(\omega) = \frac{\lambda^{|\omega|} I(\omega) \text{ is allowable w.r.t. } \alpha)}{Z_{\Delta}(\lambda)},$$
(1.2)

where allowable with respect to  $\alpha$  means that the set  $\{e \in \Delta : \omega_e = 1\} \cup \{e \in \partial \Delta : \alpha_e = 1\}$  is a matching.

It is easy to check that the monomer-dimer model satisfies the following Markov property: Let  $\sigma$  denote a random configuration on E and let  $\Delta \subset E$ . Then the conditional distribution of  $\sigma_{\Delta}$ , given  $\sigma_{E\backslash\Delta}$ , equals  $\mu_{\Lambda}^{\sigma_{\partial\Delta}}$  (and hence depends only on  $\sigma_{\partial\Delta}$ ).

The paper by van den Berg (1999) shows that the monomer-dimer model on a lattice has certain very strong spatial mixing properties. In the present paper (see Sec. III) we show explicitly how this can be used to improve, for "nice" subgraphs of a lattice, earlier results in the literature concerning the generation of (approximate) random samples. Apart from a theorem by Jerrum and Sinclair, which is stated without proof in Sec. II, and some easy to verify results on variational distance and coupling (see also Sec. II), this paper is practically self-contained.

#### II. PRELIMINARIES

In this section we give the background needed in Sec. III. First, we present some general and quite well-known results on coupling and variational distance. Then we will state the earlier mentioned result by Jerrum and Sinclair. Finally, we will present and prove a result which is very similar (but more convenient for our purpose) to a result in van den Berg.<sup>5</sup>

Throughout this section  $\Omega$  denotes an arbitrary finite set.

#### A. Coupling and variational distance

Suppose we have two probability distributions  $\mu_1$  and  $\mu_2$  on  $\Omega$ . The variational distance  $d_V(\mu_1,\mu_2)$  is defined by

$$d_{V}(\mu_{1}, \mu_{2}) = \frac{1}{2} \cdot \sum_{\omega \in \Omega} |\mu_{1}(\omega) - \mu_{2}(\omega)|. \tag{2.1}$$

Another (but equivalent) definition of variational distance is the following:

$$d_V(\mu_1, \mu_2) = \max_{A \subseteq \Omega} |\mu_1(A) - \mu_2(A)|. \tag{2.2}$$

This equivalence is quite easy to check.

Suppose we have two probability distributions  $\mu_1$  and  $\mu_2$  on  $\Omega$ . A coupling  $\mathcal{P}$  of  $\mu_1$  and  $\mu_2$  is a distribution on  $\Omega \times \Omega$  which has the following properties:

$$\sum_{\omega_1 \in \Omega} \mathcal{P}(\omega_1, \omega_2) = \mu_2(\omega_2) \text{ for all } \omega_2 \in \Omega, \tag{2.3}$$

and

$$\sum_{\omega_2 \in \Omega} \mathcal{P}(\omega_1, \omega_2) = \mu_1(\omega_1) \text{ for all } \omega_1 \in \Omega,$$
 (2.4)

i.e., the marginal distributions of  $\mathcal{P}$  are  $\mu_1$  and  $\mu_2$ . Similarly, one can define couplings of more than two probability distributions. A trivial example of a coupling is the product coupling  $\mu_1 \times \mu_2$ .

Define the event "unequal" as the set  $\{(\omega_1, \omega_2) \in \Omega \times \Omega : \omega_1 \neq \omega_2\}$ . Likewise, we define the event "equal" as the set  $\{(\omega_1, \omega_2) \in \Omega \times \Omega : \omega_1 = \omega_2\}$ . The following results, Proposition 2.1, Lemma 2.2, and Proposition 2.3, are quite standard and not difficult to prove.

Proposition 2.1: Let  $\mu_1$ ,  $\mu_2$ , and  $\mu_3$  be probability distributions on  $\Omega$ , and let  $\mathcal{P}_{1,2}$  and  $\mathcal{P}_{2,3}$  be couplings of  $\mu_1$  and  $\mu_2$ , and of  $\mu_2$  and  $\mu_3$  respectively. Then there exists a coupling  $\mathcal{P}_{1,3}$  of  $\mu_1$  and  $\mu_3$  with the following property:

$$\mathcal{P}_{1,3}(``unequal"') \leq \mathcal{P}_{1,2}(``unequal"') + \mathcal{P}_{2,3}(``unequal"'). \tag{2.5}$$

We proceed with a lemma that states some basic properties of variational distance.

Lemma 2.2: Let  $\mu$ ,  $\mu'$ , and  $\nu$  be probability distributions on  $\Omega$ . We have

- (1)  $d_V(\mu, \nu) \ge 0$ ,
- (2)  $d_V(\mu, \nu) = d_V(\nu, \mu)$ ,
- (3)  $d_V(\mu, \nu) \leq d_V(\mu, \mu') + d_V(\mu', \nu)$ ,
- (4)  $d_V(\gamma \cdot \mu + (1 \gamma) \cdot \mu', \nu) \leq \gamma \cdot d_V(\mu, \nu) + (1 \gamma) \cdot d_V(\mu', \nu)$  for all  $\gamma \in [0, 1]$ .

The following proposition relates the two notions of variational distance and couplings. Recall the notions "equal" and "unequal" defined earlier.

Proposition 2.3: For all probability distributions  $\mu$  and  $\nu$  on  $\Omega$ ,

$$d_{V}(\mu,\nu) = \min_{\mathcal{P}} \mathcal{P} (``unequal"), \tag{2.6}$$

where the minimum is taken over all couplings P of  $\mu$  and  $\nu$ .

A coupling that reaches the minimum in Proposition (2.3) is called *optimal*. For an extensive treatment of coupling methods, see Ref. 11.

#### B. Mixing times and the Jerrum-Sinclair result

Suppose we have an ergodic Markov chain on  $\Omega$ . Let  $\pi$  be the stationary distribution of this chain and let  $x \in \Omega$ . Let  $\mu^{x,t}$  denote the distribution of the Markov chain at time t, when it has started in initial state x. Let  $\epsilon > 0$ . Define the mixing time with respect to initial state x of the Markov chain as follows:

$$\tau_{x}(\epsilon) = \min_{t} \{ d_{V}(\mu^{x,t'}, \pi) \leq \epsilon \text{ for all } t' \geq t \}.$$
(2.7)

The (total) mixing time of the Markov chain is defined by

$$\tau(\epsilon) = \max_{x \in \Omega} \tau_x(\epsilon). \tag{2.8}$$

Jerrum and Sinclair<sup>9</sup> have studied the mixing time of a suitable Markov chain for the monomerdimer model. More precisely, they have proved the following: Let  $G = (\Lambda, E)$  be a finite graph, and let  $\Omega = \{\text{all matchings on } G\}$ . Consider the monomer-dimer distribution with parameter  $\lambda > 0$  on  $\Omega$ . Denote this distribution by  $\pi_{\lambda}$ . To sample from this distribution, they study a specific Markov chain  $mc_{\lambda}$  whose stationary distribution is  $\pi_{\lambda}$ . A transition  $M \to M'$  in the Markov chain  $mc_{\lambda}$  is described as follows:

(1) With probability  $\frac{1}{2}$  let M' = M; otherwise

(2) Choose uniformly at random an edge  $e = (u, v) \in E$ . Define M' as follows:

$$M' = \begin{cases} M + e & \text{if } u, v \text{ unmatched in } M, \\ M - e & \text{if } e \in M, \\ M + e - e' & \text{if either } u \text{ or } v \text{ (but noth both) is matched.} \\ & \text{and } e' \text{ is the matching edge,} \\ M & \text{otherwise} \end{cases}$$

(3) Move to M' with probability  $\min\{1, [\pi_{\lambda}(M')/\pi_{\lambda}(M)]\}$ .

Note that  $mc_{\lambda}$  is aperiodic because  $P(M,M) \ge \frac{1}{2} > 0$  for all matchings M. It is also clear that  $mc_{\lambda}$  is irreducible (because all matchings communicate via the empty matching), and easy to check that  $mc_{\lambda}$  satisfies the detailed balance condition,

$$\pi_{\lambda}(M)P(M,M') = \pi_{\lambda}(M')P(M',M). \tag{2.9}$$

We conclude that  $mc_{\lambda}$  has stationary distribution  $\pi_{\lambda}$  and that, for any initial state, the distribution of the chain converges to  $\pi_{\lambda}$ . By a clever application of the so-called canonical path method, Jerrum and Sinclair, obtained the following bound for the mixing time of  $mc_{\lambda}$ .

**Theorem 2.4:** The mixing time of  $mc_{\lambda}$  satisfies

$$\tau(\epsilon) \leq 4|E|n\lambda'(n\ln 4n + n\ln \lambda') + \ln(\epsilon^{-1})), \tag{2.10}$$

where  $\lambda' = \max\{1,\lambda\}$ , and  $n = \lfloor |\Lambda|/2 \rfloor$ .

[In fact, Proposition 12.4 of Jerrum and Sinclair,9 states

$$\tau(\epsilon) \leq 4|E|n\lambda'(n(\ln n + \ln \lambda') + \ln(\epsilon^{-1})). \tag{2.11}$$

However, we could only verify the proof when the factor  $\ln n$  is replaced by  $\ln 4n$ , in (2.11).

### C. A result on the spatial dependencies of the monomer-dimer model

The following theorem is very similar to a result in Sec. III of Ref. 5 (the ideas in which go back (Refs. 2-4), but slightly stronger and more convenient for our purpose. Therefore, and for completeness, we give a fairly detailed proof. Recall the definitions of  $V(\omega,\omega')$  and  $\deg(v)$  in Sec. I.

**Theorem 2.5:** Let, for a given value of  $\lambda$ ,  $\mu$  be the monomer-dimer distribution on a graph G = (V, E). Let  $\Delta \subset E$  and let  $\alpha, \beta \in \Omega_{\partial \Delta}$ . Then a coupling  $\mathcal{P}_{\Delta, \alpha, \beta}$  of  $\mu_{\Delta}^{\alpha}$  and  $\mu_{\Delta}^{\beta}$  exists such that

$$\mathcal{E}_{\Delta,\alpha,\beta}(|\{e \in \Delta : e \text{ edge of disagreement}\}|) \leq 2c\lambda \cdot |V(\alpha,\beta)|, \tag{2.12}$$

where  $\mathcal{E}_{\Delta,\alpha,\beta}$  denotes the expectation with respect to  $\mathcal{P}_{\Delta,\alpha,\beta}$  and c equals  $\max_{v \in \Delta} \{\deg(v)\} - 1$ . Proof: Let  $\Delta$ ,  $\alpha$ , and  $\beta$  be as in the statement of the theorem. We construct the desired coupling  $\mathcal{P}_{\Delta,\alpha,\beta}$  on  $\Omega_{\Delta} \times \Omega_{\Delta}$  as follows: Let x and y be independent configurations with distribution  $\mu_{\Delta}^{\alpha}$  and  $\mu_{\Delta}^{\beta}$  respectively. Modify these configurations in the following way. For every  $\widetilde{e} \in V(\alpha,\beta)$ , define the set

$$\widetilde{e}_{\text{diff}}(x,y) = \{e \in V(x,y): \exists a \text{ sequence } \widetilde{e} \sim e_1 \sim e_2 \sim \cdots \sim e_n = e \text{ of distinct edges in } \Delta, \text{ with } \forall i \in \{1 \cdots n\}: x_{e_i} \neq y_{e_i}.$$
 (2.13)

We call such a sequence a path of disagreement of length n from v to e, where v is the common end point of  $\tilde{e}$  and  $e_1$ . Let the set DIFF<sub> $x,y,\alpha,\beta$ </sub> be the union of paths of disagreement leaving from  $V(\alpha,\beta)$ 

$$DIFF_{x,y,\alpha\beta} := \bigcup_{\widetilde{e} \in V(\alpha,\beta)} \widetilde{e}_{diff}(x,y). \tag{2.14}$$

The modified configurations  $\tilde{x}$  and  $\tilde{y}$  are defined by

$$\widetilde{x}_e = \begin{cases} x_e & \text{if } e \in \text{DIFF}_{x,y,\alpha,\beta}, \\ y_e & \text{else}, \end{cases}$$

$$\widetilde{y}_e = y_e & \text{for all } e \in \Delta.$$
(2.15)

Note that the configurations  $\tilde{x}$  and  $\tilde{y}$  only differ from each other on DIFF<sub> $\tilde{x},\tilde{y},\alpha,\beta$ </sub> and that this set equals DIFF<sub> $x,y,\alpha,\beta$ </sub>. We define  $\mathcal{P}_{\Delta,\alpha,\beta}$  as the distribution of the pair  $(\tilde{x},\tilde{y})$  constructed as above.

Lemma 2.6: The distribution of  $\mathcal{P}_{\Delta,\alpha,\beta}$  defined above is indeed a coupling of  $\mu_{\Delta}^{\alpha}$  and  $\mu_{\Delta}^{\beta}$ .

*Proof:* Since we have defined  $\mathcal{P}_{\Delta,\alpha,\beta}$  as the distribution of  $(\tilde{x},\tilde{y})$  it is sufficient to prove that  $\tilde{x}$  has distribution  $\mu_{\Delta}^{\alpha}$  and  $\tilde{y}$  has distribution  $\mu_{\Delta}^{\beta}$ . Clearly, since  $\tilde{y}$  equals y it has distribution  $\mu_{\Delta}^{\beta}$ . It remains to show that  $\tilde{x}$  has distribution  $\mu_{\Delta}^{\alpha}$ . To do this first introduce configurations  $\hat{x}$  and  $\hat{y}$  as follows: In words,  $(\hat{x},\hat{y})$  is the pair of configurations obtained from (x,y) by exchanging x and y on the set of edges that do not have a path of disagreement to  $V(\alpha,\beta)$ . More precisely,

$$\hat{x}_e = \begin{cases} x_e & \text{if } e \in \text{DIFF}_{x,y,\alpha,\beta}, \\ y_e & \text{else}, \end{cases}$$

$$\hat{y}_e = \begin{cases} y_e & \text{if } e \in \text{DIFF}_{x,y,\alpha,\beta}, \\ x_e & \text{else}. \end{cases}$$
(2.16)

By an appropriate use of the Markov property (see the proof of Lemma 1 in Ref. 5), the pair  $(\hat{x}, \hat{y})$  has the same distribution as the pair (x, y). Finally, from the definitions of  $\hat{x}$  and  $\tilde{x}$ , it follows that  $\hat{x} = \tilde{x}$ . Hence  $\tilde{x}$  has distribution  $\mu_{\Lambda}^{\alpha}$ .

We now show that this coupling  $\mathcal{P}_{\Delta,\alpha,\beta}$  has property (2.12). First recall (see the note before Lemma 2.6) that the left-hand side of (2.12) is equal to the expected size of DIFF<sub>x,y,\alpha,\beta</sub>, where x and y are drawn independently from  $\mu_{\Delta}^{\alpha}$  and  $\mu_{\Delta}^{\beta}$ , respectively. Therefore we study the paths of disagreement for the pair (x,y). So consider an edge  $\tilde{e} \in V(\alpha,\beta)$ , say  $\tilde{e} = (v_1,v_2)$ . Observe that if a path of disagreement of length k from  $v_1$  exists, then this path is unique. (Otherwise, as one can easily check, there would be three distinct edges, which share a common endpoint, and on each of which  $x \neq y$ . But then at least two of these edges have x = 1, or at least two of these edges have y = 1, which contradicts the fact that that x and y are allowable.) For  $v_2$  a similar statement holds. Define  $l_1(\tilde{e})(l_2(\tilde{e}))$  as the path of disagreement of maximal length, starting from  $v_1$  ( $v_2$ , respectively). From the above observations we conclude that the left-hand side of (2.12) is at most

$$\sum_{\tilde{e} \in V(\alpha,\beta)} E[|l_1(\tilde{e})| + |l_2(\tilde{e})|] = \sum_{\tilde{e} \in V(\alpha,\beta)} \sum_{k=1}^{\infty} P(|l_1(\tilde{e})| \geqslant k) + P(|l_2(\tilde{e})| \geqslant k). \tag{2.17}$$

To complete the proof of Theorem 2.5 we must, in view of Eq. (2.17), bound the probability

$$P(l_1(\tilde{e}) \text{ has length} \ge k),$$
 (2.18)

and its analog for  $l_2(\tilde{e})$ . Before we do this, we first state a simple general lemma. Consider the monomer-dimer model on the very special "star-shape" graph, which consists of n edges and n+1 vertices, one of which (the "center of the star") has one edge to each of the other n vertices. It is clear that each allowable configuration has either 0 or 1 edge with value 1, and that the latter has probability  $\lambda n/(1+\lambda n)$ . Note that this is increasing in n. This observation, together with the Markov property mentioned in Sec. I (below (1.2)) implies immediately the following:

Lemma 2.7: Consider the monomer-dimer model with parameter  $\lambda$  on an arbitrary finite graph G. Let v be a vertex of G and let A be a subset of the edges of v. Then the conditional probability that there exists an edge in A with value 1, given the values of all edges outside A, is at most

$$\frac{\lambda|A|}{1+\lambda|A|}.$$

We now proceed with the proof of Theorem 2.5. Suppose a path of disagreement of length k exists. What is the conditional probability that a path of disagreement of length k+1 exists? Let  $e_1 \sim \cdots \sim e_k = e$  be the (unique) path of length k leaving from  $v_1$ , so that  $\tilde{e} \sim e_1$ . By the uniqueness property mentioned before, we have that the path of disagreement of length k+1 (if it exists) is an extension of the path of length k. Define

$$Adi(e) = \{b \in \Delta : b \sim e \text{ and } b \neq e_{k-1}\}. \tag{2.19}$$

Note that  $|\operatorname{Adj}(e)| \le c$ , with c as in the statement of the theorem.

By assumption  $x_e = 0$  and  $y_e = 1$  or vice versa. Without loss of generality we assume the former. Since y is a matching, we have  $y_b = 0$  for every edge  $b \in \mathrm{Adj}(e)$ . Hence we have a path of disagreement of length k+1, if and only if an edge  $a \in \mathrm{Adj}(e)$  exists with  $x_a = 1$ . By Lemma 2.7 above, the (conditional) probability of this event is at most  $\lambda c/\lambda c + 1$ . Iterating the above we get

$$\mathcal{P}_{\Delta,\alpha,\beta}(l_1(\tilde{e})\text{has length} \ge k) \le \left(\frac{\lambda c}{\lambda c + 1}\right)^k.$$
 (2.20)

The same result holds for  $l_2(\tilde{e})$ .

Combining (2.20) with (2.17), it follows that the left-hand side of (2.12) is at most

$$2\sum_{\tilde{c} \in V(\alpha, \beta)} \sum_{k=1}^{\infty} \left( \frac{\lambda c}{\lambda c + 1} \right)^{k} = 2\lambda c \cdot |V(\alpha, \beta)|.$$
 (2.21)

This completes the proof of Theorem 2.5.

Remark 2.8: In Sec. III A and III B, we will only work with d-dimensional hypercubes  $\Delta$ . For such sets  $\Delta$ , each edge on the boundary  $\partial \Delta$  has exactly one vertex in common with an edge in the box  $\Delta$ . For these special cases, the above result is improved by a factor 2, so that

$$\mathcal{E}_{\Delta,\alpha,\beta}(\#\{e \in \Delta : e \text{ edge of disagreement}\}) \leq \lambda(2d-1) \cdot |V(\alpha,\beta)|, \tag{2.22}$$

for every hypercube  $\Delta$ .

#### III. RANDOM SAMPLING ON SUBGRAPHS OF THE d-DIMENSIONAL LATTICE

#### A. Description and motivation of the method

In Sec. III B we stated the Jerrum–Sinclair result. This result holds for general graphs. In the present section we study certain specifically "nice" graphs, say a d-dimensional torus (described more precisely below). Suppose we want to sample (approximately) from the monomer–dimer model for such a graph. According to the Jerrum–Sinclair result (Theorem 2.4) we can do this by running the Markov chain  $mc_{\lambda}$  (described in Sec. II B) a number of steps given by (2.11). For the torus this is, for fixed  $\lambda$  and  $\epsilon$ , asymptotically of order (Volume)<sup>3</sup>×log(Volume). Here Volume is the number of edges in the graph (or the number of vertices, which for these graphs differs a constant factor from the number of edges).

Can this, for these special graphs, be improved? There are several possibilities. One approach is to use logarithmic Sobolev inequalities: the results on spatial dependencies in Sec. II C imply a

mixing condition which, in turn, following a quite general theory developed by Stroock and Zegarlinski (see Ref. 15), could lead to a bound on the mixing time of order Volume ×log(Volume). (We write *could* because there is an extra, quite subtle, condition which has to be checked to obtain such a bound from the Stroock–Zegarlinski theory; see Theorem 1 in the survey paper, by Frigessi, Martinelli, and Stander.) This result would be very interesting, but when one really wants to generate random samples, one not only wants to know the *asymptotic order* of the mixing time, but one needs an *explicit upper bound* to carry out the algorithm. To get (reasonable) explicit bounds from the Stroock–Zegarlinski theory is probably a lot of work which (in our opinion) is certainly worth the effort.

However, in the present paper we follow a somewhat different approach, which is based on a small modification of coupling and rescaling arguments which have become quite standard (see Aizenman and Holley, and Martinelli and Olivieri. This approach has the advantage that it gives, with relatively simple and few computations, an explicit bound whose asymptotic order is 'only a little worse' than the above mentioned Volume×log(Volume). (We get an extra factor of order log(Volume).)

Our approach is to combine (using rescaling and coupling arguments) the result of Jerrum and Sinclair (Theorem 2.4) with the result on spatial dependencies in Sec. II C. Although this approach applies to a larger class of graphs (see Remark 2 Sec. III C), we concentrate for simplicity on a graph  $\Gamma$ , which corresponds to a d-dimensional torus. More precisely, let N be a positive integer, and define  $\Gamma$  as the pair  $(\Lambda_{\Gamma}, \mathcal{E}_{\Gamma})$  where the set of vertices  $\Lambda_{\Gamma}$  is defined as

$$\Lambda_{\Gamma} := \{0, \dots, N\}^d, \tag{3.1}$$

and the set of edges  $E_{\Gamma}$  is

$$E_{\Gamma} := \{ (v_1, v_2) : v_1, v_2 \in \Lambda_{\Gamma} \text{ and } |v_1 - v_2| = 1 \pmod{(N-1)} \}, \tag{3.2}$$

where  $|\cdot|$  denotes the  $l_1$  distance. We would like to sample from the monomer-dimer distribution  $\pi_{\Gamma}$  with parameter  $\lambda$  on this graph.

One way of approximate sampling from this distribution on  $\Gamma$  is the following: Let  $\delta > 0$ . Let  $\Delta$  be a d-dimensional cube of length l. (Here l depends on  $\lambda$  and d; a suitable value will be determined later.) More precisely,  $\Delta$  is the following set of edges:

$$\Delta := \{ (v_1, v_2) : v_1, v_2 \in \{0, \dots, l_n\}^d \text{ and } |v_1 - v_2| = 1 \}.$$
 (3.3)

Let  $\hat{X}(t)$ , t = 0,1,... be the Markov chain with state space  $\{0,1\}^{E_{\Gamma}}$ , which starts in some  $x_0 \in \Gamma$ , and of which the transitions are described as follows: Suppose  $\hat{X}(t) = x$ . Choose u.a.r. a vertex  $i \in \Lambda_{\Gamma}$ . Let  $\widetilde{\Delta}$  be the box  $\Delta$  shifted over i in the torus, i.e.,

$$\Delta = \{ ((v_1 + i) \pmod{N}, (v_2 + i) \pmod{N}) : (v_1, v_2) \in \Delta \}.$$
(3.4)

Consider the monomer-dimer distribution on  $\widetilde{\Delta}$  with boundary condition  $x_{\partial\widetilde{\Delta}}$  (and parameter  $\lambda$ ), denoted by  $\mu_{\widetilde{\Delta}}^{x_{\partial\widetilde{\Delta}}}$ . Now sample a configuration  $\widetilde{x}$  from this distribution. At time t+1 the state becomes

$$\hat{X}(t+1)_e = \begin{cases} \hat{X}(t)_e & \text{if } e \in \widetilde{\Delta}, \\ \widetilde{x}_e & \text{if } e \in \widetilde{\Delta}. \end{cases}$$
(3.5)

It can be proved, using the spatial mixing properties mentioned before, that for l sufficiently large, the mixing time of this Markov chain for fixed  $\lambda$  is of order  $\mathcal{O}(|\Lambda_{\Gamma}| \cdot \log |\Lambda_{\Gamma}|)$ , i.e., of the same order we mentioned above in connection with logarithmic Sobolev inequalities.

However, a problem arises when one actually tries to execute this algorithm. How to compute the above mentioned distribution  $\mu_{\tilde{\Delta}}^{x_{\tilde{\sigma}\tilde{\Delta}}}$ ? Even for relatively small  $\Delta$ , this is a huge problem. For example, if d=2 and the length l of the hypercube is 10, the state space of  $\mu_{\tilde{\Delta}}^{x_{\tilde{\sigma}\tilde{\Delta}}}$  has already more than  $2^{100}$  elements. In practice, this algorithm cannot be used.

One way to proceed now would be to use certain comparison theorems to obtain a bound on the mixing time of the Markov chain  $mc_{\lambda}$  for this model from the bound on the mixing time of the above "block dynamics" (see Diaconis and Saloff-Coste, Randall and Tetali, <sup>14</sup> and Martinelli<sup>12</sup>). However, these comparison arguments do not involve the two mixing times directly but indirectly (via the spectral gap or logarithmic Sobolev constant). Since the relation between the mixing time (2.7)-(2.8) and these quantities is not tight, this method would introduce a factor of order Volume, so the final result would be of order (Volume)  $^2 \log(\text{Volume})$ .

Therefore we do the following: Instead of drawing a configuration *exactly* from the distribution  $\mu_{\widetilde{\Delta}}^{x_{\partial \widetilde{\Delta}}}$  mentioned before, we will sample *approximately* from this distribution. In other words, we replace each (*macro*) step in the Markov chain by a number of *micro* steps where each micro step corresponds with a transition of the Markov chain  $mc_{\lambda}$  (on  $\widetilde{\Delta}$ , with boundary condition  $x_{\partial \widetilde{\Delta}}$ ) studied by Jerrum and Sinclair. It will turn out that (for fixed  $\delta$  and  $\lambda$ ), the total number of micro steps needed to obtain a " $\delta$ -close" approximate sample from  $\pi_{\Gamma}$  is at most of order Volume (log(Volume))<sup>2</sup> (see Corollary 3.3 at the end of this section).

More precisely, the modified Markov chain, which we denote by X(t), t=0,1,..., has the same state space and initial state as  $\hat{X}(t)$ , but the transitions are now as follows: Suppose X(t)=x. As before, choose u.a.r. a vertex  $i \in \Gamma$ ; determine the box  $\tilde{\Delta}$ , and consider the monomer-dimer distribution  $\mu_{\tilde{\Lambda}}^{x_{\partial}\tilde{\Delta}}$ . We will approximate this distribution. To do this, first define

$$\epsilon = \left(\frac{|\Delta| - \lambda(2d - 1)|\partial\Delta|}{|E_{\Gamma}||\Delta|}\right) \frac{\delta}{2}.$$
(3.6)

The choice of this value will become clear later. Now consider the (auxiliary) Markov chain  $mc_{\lambda}$  (with respect to the monomer-dimer model on  $\widetilde{\Delta}$ , with boundary condition  $x_{\partial\widetilde{\Delta}}$ ) described in Sec. II B. Although the initial state does not matter in the computations below, it is natural to take it equal to  $x_{\widetilde{\Delta}}$ . Denote the distribution of this chain at time t by  $\nu_{\widetilde{\Delta}}^{t,x}$ . Let  $\Delta^*$  be the set of vertices which are endpoints of edges of  $\Delta$ . From Theorem 2.4 it follows that  $\nu_{\widetilde{\Delta}}^{t,x}$  converges to  $\mu_{\widetilde{\Delta}}^{x_{\partial\widetilde{\Delta}}}$ , and that, if the number of steps made by that Markov chain is at least T, given by

$$T := |\Delta| |\Delta^*| \lambda' [|\Delta^*| \ln(2|\Delta^*|) + (|\Delta^*|) \ln \lambda' + 2 \ln(\epsilon^{-1})], \tag{3.7}$$

then

$$d_{V}(\nu_{\widetilde{\Delta}}^{T,x},\mu_{\widetilde{\Delta}}^{x_{\partial\widetilde{\Delta}}}) \leq \epsilon. \tag{3.8}$$

Let  $\tilde{x}$  be the configuration on  $\tilde{\Delta}$  after T transitions; this is a sample from  $\nu_{\tilde{\Delta}}^{T,x}$ . Now take

$$X(t+1)_e = \begin{cases} X(t)_e & \text{if } e \in \widetilde{\Delta} \\ \widetilde{x}_e & \text{if } e \in \widetilde{\Delta} \end{cases}$$
 (3.9)

This completes our description of a (macro) step in the Markov chain X.

In the next section we will give an upper bound for the number of macro steps after which the variational distance between  $\mu^t$  (the distribution of X(t) at time t) and  $\pi_{\Gamma}$  becomes smaller than  $\delta$ . The total number of micro steps needed then simply follows from multiplying this by the number T in (3.7).

## B. A bound on the number of steps

In this section we will bound the number of steps of the Markov chain X(t) to approximately reach the stationary distribution  $\pi_{\Gamma}$ . First, we define a suitably coupled system (X(t),Y(t)), t=0,1,..., where X(t) is the Markov chain introduced in the previous subsection, and Y(t) is a Markov chain with the same transition probabilities as  $\hat{X}$  (which was also introduced in the previous subsection), but which starts with the stationary distribution  $\pi_{\Gamma}$  (and hence keeps this distribution). Using the results in Sec. II, we will obtain an upper bound for the variational distance between the distributions of X(t) and Y(t) for every time t. This is done by studying the number of edges of disagreement |V(X(t),Y(t))|.

More precisely, let  $X(0)=x_0$  and let Y(0) be drawn from the distribution  $\pi_{\Gamma}$ . Suppose at time t,X(t)=x and Y(t)=y. Now we follow the description of a transition of X(t) given in the previous subsection. However, instead of sampling a single configuration  $\tilde{x}$  on  $\tilde{\Delta}$ , we now sample a pair  $(\tilde{x},\tilde{y})$  as follows. First consider the following three distributions on  $\Omega_{\tilde{\Delta}}$ :  $\nu_{\tilde{\Delta}}^{T,x}$ ,  $\mu_{\tilde{\Delta}}^{x_0\tilde{\Delta}}$  and  $\mu_{\tilde{\lambda}}^{y_0\tilde{\Delta}}$ . Let

$$\mathcal{P}_{\text{opt}\nu_{\tilde{\Lambda}}^{T,x},\mu_{\tilde{\Lambda}}^{x,\tilde{\sigma}_{\tilde{\Lambda}}}} \tag{3.10}$$

be an optimal coupling of  $\nu_{\widetilde{\Delta}}^{T,x}$  and  $\mu_{\widetilde{\Delta}}^{x_{\partial\widetilde{\Delta}}}$ , and

$$\mathcal{P}_{\mu_{\tilde{\Lambda}}^{x_{\tilde{\sigma}\tilde{\Lambda}}}, \mu_{\tilde{\Lambda}}^{y_{\tilde{\sigma}\tilde{\Lambda}}}}$$
 (3.11)

be a coupling of  $\mu_{\tilde{\Lambda}}^{x_{\tilde{\sigma}}\tilde{\Delta}}$  and  $\mu_{\tilde{\Lambda}}^{y_{\tilde{\sigma}}\tilde{\Delta}}$  which satisfies Theorem 2.5. Finally, let

$$\mathcal{P}_{\nu_{\tilde{\Delta}}^{T,x},\mu_{\tilde{\Delta}}^{Y_{\partial\tilde{\Delta}}}} \tag{3.12}$$

be a coupling of  $\nu_{\widetilde{\Delta}}^{T,x}$  and  $\mu_{\widetilde{\Delta}}^{y_{\partial\widetilde{\Delta}}}$  obtained from the two previous couplings as described in the proof of Proposition 2.1. The expectation with respect to the distribution (3.10) is denoted by  $\mathcal{E}_{\text{opt},\nu_{\widetilde{\Delta}}^{T,x},\mu_{\widetilde{\Delta}}^{x_{\partial\widetilde{\Delta}}}}$ . The expectations for the other two couplings are denoted similarly. Now, sample a pair  $(\widetilde{x},\widetilde{y})$  from this last coupling (3.12). Now take

$$X(t+1)_{e} = \begin{cases} X(t)_{e} & \text{if } e \in \widetilde{\Delta}, \\ \widetilde{x}_{e} & \text{if } e \in \widetilde{\Delta}. \end{cases} \quad Y(t+1)_{e} = \begin{cases} Y(t)_{e} & \text{if } e \in \widetilde{\Delta}, \\ \widetilde{y}_{e} & \text{if } e \in \widetilde{\Delta}. \end{cases}$$
(3.13)

This completes the description of the transitions of the pair (X(t), Y(t)). Note that  $\widetilde{X}$  has been drawn from  $\nu_{\widetilde{\Delta}}^{T,x}$  so that the Markov chain X(t) has indeed the same transition probabilities as in Sec. III A. Similarly, note that Y(t) has indeed distribution  $\pi_{\Gamma}$  for each t.

Let  $\mu^t$  denote the distribution of X(t). Let E(t) denote the expectation of |V(X(t),Y(t))|. Using Proposition 2.3 we have

$$d_{V}(\mu^{t}, \pi_{\Gamma}) \leq P[X(t) \neq Y(t)] \leq E(t). \tag{3.14}$$

Therefore we will study E(t). In particular, we study the change in this quantity after one (macro) step of the coupled Markov chain.

Using a property analogous to Eq. (2.5) in Proposition 2.1 we get

$$\begin{split} \mathcal{E}_{\nu_{\widetilde{\Delta}}^{T,x},\mu_{\widetilde{\Delta}}^{y_{\delta\widetilde{\Delta}}}}(|\{e\in\widetilde{\Delta}:e \text{ edge of disagreement}\}|) \\ &\leq \mathcal{E}_{\mathrm{opt},\nu_{\widetilde{\Delta}}^{T,x},\mu_{\widetilde{\Delta}}^{x_{\delta\widetilde{\Delta}}}}(|\{e\in\widetilde{\Delta}:e \text{ edge of disagreement}}\}|) \\ &+ \mathcal{E}_{\mu_{\widetilde{\Delta}}^{x_{\delta\widetilde{\Delta}}},\mu_{\widetilde{\Delta}}^{y_{\delta\widetilde{\Delta}}}}(|\{e\in\widetilde{\Delta}:e \text{ edge of disagreement}}\}|). \end{split} \tag{3.15}$$

So we need upper bounds for the expectations in the right-hand side of (3.15). By Theorem 2.5 (and Eq. (2.22)),

$$\mathcal{E}_{\mu_{\widetilde{\Lambda}}^{x,\partial\widetilde{\Delta}},\mu_{\widetilde{\Lambda}}^{y,\widetilde{a}\widetilde{\Delta}}}(\big|\big\{e\in\widetilde{\Delta}:e\ \text{ edge of disagreement}\big\}\big|) \leq \lambda(2d-1)\big|V(x_{\partial\widetilde{\Delta}},y_{\partial\widetilde{\Delta}})\big|. \tag{3.16}$$

Because the coupling  $\mathcal{P}_{\mathrm{opt},\, \nu_{\widetilde{\Delta}}^{T,x},\, \mu_{\widetilde{\lambda}}^{x_{\varrho\widetilde{\Delta}}}}$  is optimal, we have

$$\begin{split} \mathcal{E}_{\mathrm{opt},\,\nu_{\widetilde{\Delta}}^{T,x},\,\mu_{\widetilde{\Delta}}^{x_{\vartheta\widetilde{\Delta}}}(|\{e\in\widetilde{\Delta}:e\ \text{ edge of disagreement}\}|) & \leqslant |\widetilde{\Delta}|\cdot\mathcal{P}_{\mathrm{opt},\,\nu_{\widetilde{\Delta}}^{T,x},\,\mu_{\widetilde{\Delta}}^{x_{\vartheta\widetilde{\Delta}}}(\text{``unequal''})} \\ & = |\widetilde{\Delta}|\cdot d_{V}(\,\nu_{\widetilde{\Delta}}^{T,x},\,\mu_{\widetilde{\Delta}}^{x_{\vartheta\widetilde{\Delta}}}) \\ & \leqslant \epsilon\cdot|\widetilde{\Delta}|\,. \end{split}$$

The last inequality follows from (3.8). Together, Eqs. (3.15)-(3.17) yield

$$\mathcal{E}_{\nu_{\widetilde{\Delta}}^{T,x},\mu_{\widetilde{\Delta}}^{y,\delta}}(\big|\big\{e\in\widetilde{\Delta}:e\ \text{ edge of disagreement}\big\}\big|)\leqslant\epsilon\cdot\big|\widetilde{\Delta}\big|+\lambda\cdot(2d-1)\big|V(x_{\partial\widetilde{\Delta}},y_{\partial\widetilde{\Delta}})\big|. \quad (3.18)$$

We now state and prove the following Lemma:

Lemma 3.1:

$$E(t+1) \le b \cdot E(t) + \epsilon |\Delta|, \tag{3.19}$$

where

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$$b \coloneqq 1 - \frac{|\Delta| - \lambda(2d - 1)|\partial\Delta|}{|E_{\Gamma}|}.$$
 (3.20)

*Proof:* Let M(t) = |V(X(t), Y(t))|. Note that the expectation of M(t) is equal to E(t). Suppose that X(t), Y(t), i and hence  $\widetilde{\Delta}$  are known. Consider the conditional expectation of the number of edges of disagreement that disappear during the transition  $t \rightarrow t+1$ ,

$$E[M(t) - M(t+1)|X(t) = x, Y(t) = y, \widetilde{\Delta} = A] = |V(x_A, y_A)| - \mathcal{E}_{\nu_A^{T, x}, \mu_A^{y, y_A}}(|V(\widetilde{x}, \widetilde{y})|). \quad (3.21)$$

By (3.18) this is larger than or equal to

$$|V(x_A, y_A)| - \epsilon |A| - \lambda (2d - 1)|V(x_{\partial A}, y_{\partial A})|. \tag{3.22}$$

Averaging over A we get

$$E[M(t)-M(t+1)|X(t)=x,Y(t)=y] \ge E(|V(x_{\tilde{\Delta}},y_{\tilde{\Delta}})|) - \epsilon|\Delta| - \lambda(2d-1)E(|V(x_{\partial\tilde{\Delta}},y_{\partial\tilde{\Delta}})|), \tag{3.23}$$

where the expectation in the right-hand side refers to the distribution of  $\tilde{\Delta}$ . Recall that this is the uniform distribution, so that by symmetry each edge of  $E_{\Gamma}$  has the same probability,  $|\Delta|/|E_{\Gamma}|$ , to belong to  $\tilde{\Delta}$ . Similarly, the probability that a given edge belongs to  $\partial \tilde{\Delta}$  equals  $|\partial \Delta|/|E_{\Gamma}|$ . Hence the right-hand side of (3.23) equals

$$\frac{|\Delta|}{|E_{\Gamma}|} \cdot |V(x,y)| - \epsilon |\Delta| - \lambda (2d - 1) \frac{|\partial \Delta|}{|E_{\Gamma}|} |V(x,y)|.$$

So we have

$$E[M(t) - M(t+1)|X(t), Y(t)] \ge M(t) \left( \frac{|\widetilde{\Delta}| - \lambda(2d-1)|\partial \Delta|}{|E_{\Gamma}|} \right) - \epsilon |\Delta|. \tag{3.24}$$

Taking expectations in (3.24), we get

$$E(t) - E(t+1) \ge E(t) \left( \frac{|\Delta| - \lambda(2d-1)|\partial\Delta|}{|E_{\Gamma}|} \right) - \epsilon |\Delta|, \tag{3.25}$$

from which the lemma follows immediately.

For the moment we assume that the following inequalities:

$$0 < 1 - \frac{|\Delta| - \lambda(2d - 1)|\partial\Delta|}{|E_{\Gamma}|} < 1 \tag{3.26}$$

hold, and will come back to this later.

Iterating Eq. (3.19) we get

$$E(t+1) \leq b^{t+1} |E_{\Gamma}| + \epsilon |\Delta| \cdot \sum_{i=0}^{t} b^{i} = b^{t+1} |E_{\Gamma}| + \frac{(1-b^{t+1})}{1-b} \epsilon |\Delta| = b^{t+1} |E_{\Gamma}| + (1-b^{t+1}) \frac{\delta}{2}.$$
(3.27)

Here we use the definition of  $\epsilon$  and b (see Eqs. (3.6) and (3.20)), and the fact that  $E(0) \leq |E_{\Gamma}|$ . With (3.14) this gives

$$d_V(\mu^t, \pi_\Gamma) \le b^t |E_\Gamma| + (1 - b^t) \frac{\delta}{2}.$$
 (3.28)

If we want to find t, such that the above mentioned variational distance is smaller than  $\delta$ , it suffices to solve

$$b^t|E_{\Gamma}| \leqslant \frac{\delta}{2}.\tag{3.29}$$

Taking logarithms on both sides of (3.29) and using that  $ln(1-x) \le -x$  for 0 < x < 1, we find that (3.29) holds if

$$t \ge \frac{(\ln(2|E_{\Gamma}|) + \ln(\delta^{-1})) \cdot |E_{\Gamma}|}{|\Delta| - \lambda(2d - 1)|\partial\Delta|}.$$
(3.30)

Recall that every step of the Markov chain X(t) is in fact a macro step which corresponds with T micro steps in some box  $\widetilde{\Delta}$ , where T is given by (3.7). Hence the total number of micro steps  $\tau(\delta)$  after which the distribution of X(t) has variational distance  $\leq \delta$  from  $\pi_{\Gamma}$  is at most T times the right-hand side of (3.30), i.e.,

$$\tau(\delta) \leq |\Delta| |\Delta^*| \lambda' \cdot [|\Delta^*| \ln(2|\Delta^*|) + |\Delta^*| \ln(\lambda') + 2\ln(\epsilon^{-1})] \times \left( \frac{(\ln(2|E_{\Gamma}|) + \ln(\delta^{-1}))|E_{\Gamma}|}{|\Delta| - \lambda(2d-1)|\partial\Delta|} \right), \tag{3.31}$$

where  $\epsilon = \epsilon(\delta)$  is defined as in (3.6). Optimization considerations on a simplified modification of the right-hand side of (3.31) lead to the following choice of the length l of  $\Delta$ :

$$l \coloneqq [\lambda(4d+2)]. \tag{3.32}$$

Note that  $|\Delta^*| = (l+1)^d$ ,  $|\Delta| = dl(l+1)^{d-1}$ , and  $|\partial\Delta| = 2d(l+1)^{d-1}$ , so that, with l given by (3.32),

$$|\Delta^*| = ([\lambda(4d+2)]+1)^d, \tag{3.33}$$

$$|\Delta| = d[\lambda(4d+2)]([\lambda(4d+2)]+1)^{d-1},$$
 (3.34)

$$|\partial \Delta| = 2d([\lambda(4d+2)]+1)^{d-1}.$$
 (3.35)

Using (3.33)–(3.35), it is easy to check that, for every  $\lambda > 0$  and every  $d \ge 2$ , the above choice of l implies the upper bound in (3.26). The lower bound in (3.26) is satisfied if  $|E_{\Gamma}| > |\Delta|$ , i.e. (in terms of  $\lambda$  and d) if  $|E_{\Gamma}|$  is larger than the right-hand side of (3.34). Using (3.33), (3.34), (3.35), and (3.6), we can now express the upper bound (3.31) on  $\tau(\delta)$  completely in terms of  $\delta$ ,  $\lambda$ , d, and  $|E_{\Gamma}|$ .

Summary of the algorithm and the main result. Concluding, we can state the following: Let  $0 < \delta < 1$  and  $\lambda > 0$ . Consider the monomer-dimer distribution  $\pi_{\Gamma}$  with parameter  $\lambda$  on the d-dimensional torus  $\Gamma$ , as described in Sec. III A. Take  $l = \lceil (4d+2)\lambda \rceil$  and let  $\Delta$  be the hypercube of length l as described in Sec. III A. Compute  $\epsilon$  from (3.6). Finally, compute T for the above choice of l, as in (3.7). Consider the Markov chain X(t) with state space  $\{0,1\}^{E_{\Gamma}}$ , with transitions described as follows: Choose u.a.r. a vertex  $i \in \Gamma$  and consider the box  $\widetilde{\Delta} = (i + \Delta)$ . On this box (with the current X(t) values on the boundary fixed) run the Markov chain  $mc_{\lambda}$  (described in Sec. II B) for T steps. These steps are called micro steps. This completes one transition (macro step) in the Markov chain X(t).

**Theorem 3.2:** In the algorithm described above, the number of micro steps  $\tau(\delta)$  after which the distribution of X(t) has variational distance smaller than or equal to  $\delta$  from the stationary distribution  $\pi_{\Gamma}$  satisfies

$$\tau(\delta) \leq T \cdot \frac{(\ln(2|E_{\Gamma}|) + \ln(\delta^{-1}))|E_{\Gamma}|}{|\Delta| - \lambda(2d - 1)|\partial\Delta|}$$

$$= |\Delta||\Delta^*|\lambda' \cdot [|\Delta^*|\ln(2|\Delta^*|) + |\Delta^*|\ln(\lambda') + 2\ln(\epsilon^{-1})]$$

$$\times \left(\frac{(\ln(2|E_{\Gamma}|) + \ln(\delta^{-1}))|E_{\Gamma}|}{|\Delta| - \lambda(2d - 1)|\partial\Delta|}\right), \tag{3.36}$$

where  $|\Delta^*|$ ,  $|\Delta|$ ,  $|\partial\Delta|$  and  $\epsilon$  are given by (3.33), (3.34), (3.35), and (3.6), respectively, and  $\lambda' = \max(\lambda, 1)$ .

This result gives immediately (note the dependence of  $\epsilon$  on  $|E_{\Gamma}|$ ) Corollary 3.3: For the algorithm above, if  $\lambda$ , d and  $\delta$  are fixed,  $\tau(\delta)$  satisfies

$$\tau(\delta) = \mathcal{O}(|E_{\Gamma}|(\ln(|E_{\Gamma}|))^2). \tag{3.37}$$

Remark 3.4: From (3.37) it follows that, for fixed  $\lambda$  and  $\delta$ , on a large torus our bound is considerably better than the bound of Jerrum and Sinclair (Theorem 2.4). (Note that on a torus, the number of edges equals the dimension times the number of vertices, so  $|E_{\Gamma}| = d|\Lambda_{\Gamma}|$ .) However,

our bound (3.31) involves a factor  $\lambda^{2d}\lambda'$ , while the bound of Jerrum and Sinclair is linear in  $\lambda'$ , which is important for certain applications (see Ref. 9). Hence if the size of the torus is relatively small with respect to  $\lambda$ , their bound is better than ours.

#### C. Remarks

- (1) Since the definition of the Markov chain X(t) depends on  $\delta$ , it is, strictly speaking, not correct to call  $\tau(\delta)$  in (3.31) its mixing time.
- (2) The algorithm in the previous section, was described for a torus  $\Gamma$ . A similar result is still valid when the algorithm is applied to a sufficiently nice finite subset of  $\mathbf{Z}^d$ , for instance a hypercube  $\Gamma = (\Lambda_{\Gamma}, E_{\Gamma})$ , where  $\Lambda_{\Gamma} = \{0,...,m\}^d$  and  $E_{\Gamma} = \{(v_1,v_2):v_1,v_2 \in \Lambda_{\Gamma} \text{ and } | v_1-v_2 | = 1\}$ . Since  $\Gamma$  is not a torus, the box  $\widetilde{\Delta}$  must now be defined as  $\widetilde{\Delta} = (i+\Delta) \cap E_{\Gamma}$ , where the vertex i is now the center of the box  $\widetilde{\Delta}$ . The fact that in some cases  $\widetilde{\Delta}$  consists of roughly  $|\Delta|/2^d$  elements leads to an increase of the size of a suitable  $\Delta$ . This in turn leads to a number of micro steps needed in the procedure which is a constant (depending on the dimension d) larger than that for our torus.
- (3) One may think of several modifications of our computations to improve (decrease) the right-hand side of (3.36). For instance it would be interesting and worth trying to improve Theorem 2.5. As to alternative methods, see the remark about logarithmic Sobolev inequalities in the beginning of this Section.

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